

# An Execution Analyzer for the Sigma 5 Computer

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*Since many different computers on the market today claim to perform the same class of tasks, and no uniform criterion has been established to aid the decision-making process, the problem of how to select the "best" computer for a particular job mix becomes almost a matter of personal preference and chance. One technique for evaluating the performance of a central processing unit is based on the frequency of usage of the various machine instructions. This technique is particularly applicable to machines used in a dedicated process-control activity such as those of the DSN Tracking Stations. This article describes a program which has been written for the Sigma 5 computer to gather data on the dynamic usage of computer instructions in various tasks.*

## I. Introduction

The technique for evaluating the performance of a central processing unit based on the frequency of usage of the various machine instructions is particularly applicable to machines used in a dedicated process-control activity such as those of the DSN Tracking Stations. This article describes a program written for the Sigma 5 computer to gather data on the dynamic usage of computer instructions in various tasks. Results of a study using this program are described in a companion article by Layland and Klimasauskas<sup>1</sup>.

This data gathering program is named MR $\phi\phi$ 7, and its overall operation is described in Section II. Subsequent sections include instructions for MR $\phi\phi$ 7 operation, and for attaching it to various types of user programs; a description of the format in which the data gathered are written into a disk file; and a description of error conditions which may be encountered, and their associated messages. Section V describes in some detail the set of

programs which have been used to transform the data gathered by MR $\phi\phi$ 7 into a visually meaningful format, and presents an example of that output format.

## II. The Interpreter

MR $\phi\phi$ 7 is an interpretive executor which dynamically tabulates instruction usage. The entire problem-solving repertoire of Sigma 5 instructions is allowed.<sup>2</sup> Results from program execution are stored in a disk file for further analysis. Instructions which communicate with the monitor (CALs) are limited to background mode only. Monitor CALs dealing with interrupts, program control transfers, and traps are excluded.

MR $\phi\phi$ 7 may be tailored to operate in either a Batch Processing Monitor (BPM) or a Batch Timesharing Monitor (BTM) environment through an assembly control parameter. This parameter (P:USER), when set equal to :BPM, produces code to interpret BPM monitor CALs.

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<sup>1</sup>Layland, J. W., and Klimasauskas, C. C., *A Myopic View of Computer-Based System Design* (to be published).

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<sup>2</sup>These do not include control, and input/output instructions.

When set equal to :BTM, it produces code to interpret BTM monitor CALs.

Since this program was designed as a tool to study a wide range of problems, certain features (described in Section VI) were not considered of sufficient importance to warrant implementation. If these or other features (such as foreground CALs) must be used, their code may be bypassed through the use of a HOLD-GO pair (see Section III).

To facilitate analysis of data from MR $\phi\phi$ 7, a series of programs have been written which transform the data into informative formats. These are described in Section V.

### III. Operating Instructions

#### A. FORTRAN Programs

MR $\phi\phi$ 7 is easily attached to any extended FORTRAN IV program. To activate it, simply use this statement:

```
CALL MR $\phi\phi$ 7
:
STOP
```

If it is desired to deactivate MR $\phi\phi$ 7 with intentions of restarting it, use the following sequence:

```
CALL MR $\phi\phi$ 7
*   } Program executed under control of
*   } MR $\phi\phi$ 7
*
CALL HOLD
:   } Program executed without control of
:   } MR $\phi\phi$ 7
CALL GO
*   } Program executed under control of
*   } MR $\phi\phi$ 7
STOP
```

MR $\phi\phi$ 7 also provides the option of writing its current data tables to M:EO, clearing them, and relinquishing control of the program. This effectively provides the user the capability of selectively looking at the instruction repertoire of several very different portions of the program. A word of caution: MR $\phi\phi$ 7 *must* be in control of the program when it exits to the monitor. If MR $\phi\phi$ 7 is not in control when the program is to be terminated, the data file will be lost. To close and save the output file, use "CALL WDATCLSE." This routine will exit to the

monitor on closing the file. Note that all user output files which are open will be lost. A sample sequence is:

```
CALL MR $\phi\phi$ 7
*   } Program executed under MR $\phi\phi$ 7
*   } control
*
CALL MR $\phi\phi$ 7
:   } Data tables written to
:   } M:EO ("DEXEC"), cleared, and
:   } control relinquished
:   }
:   } Program executed without control
:   } of MR $\phi\phi$ 7
CALL MR $\phi\phi$ 7
*   } Program executed under MR $\phi\phi$ 7
*   } control
*
CALL MR $\phi\phi$ 7
:   } Data tables written to
:   } M:EO ("DEXEC"), cleared, and
:   } control relinquished
:   }
:   }
CALL MR $\phi\phi$ 7
STOP   } End sequence giving MR $\phi\phi$ 7
        } control
```

It is allowable to use any number of HOLD-GO pairs within a MR $\phi\phi$ 7 pair.

#### B. Assembly Programs

MR $\phi\phi$ 7 is attached to assembly programs using the assembly language equivalent of the FORTRAN calling sequences above, namely:

```
CALL MR $\phi\phi$ 7
is replaced by
REF MR $\phi\phi$ 7
LI,14 0
BAL,15 MR $\phi\phi$ 7
```

```
CALL GO
is replaced by
REF GO
LI,14 0
BAL,15 GO
```

```
CALL HOLD
is replaced by
REF HOLD
LI,14 0
BAL,15 HOLD
```

```
STOP
is replaced by
M:EXIT
```

## IV. Data Record Format

MR $\phi\phi$ 7 writes the data tables to a disk file ("DEXEC") through the M:EO DCB. The DCB does not need to be assigned through an ASSIGN card, nor should it be. For each MR $\phi\phi$ 7 pair and for the MR $\phi\phi$ 7-STOP pair, a data record is written. If data already exist in file "DEXEC," they will be lost. The format of each record is as shown in Fig. 1.

## V. Support Programs

To facilitate the analysis of data from MR $\phi\phi$ 7, a number of support programs have been written to transform the data to usable and meaningful terms. These programs are briefly described below.

### A. PUNCH

PUNCH is a FORTRAN program which takes the file DEXEC produced by MR $\phi\phi$ 7 and generates records in the standard binary format acceptable to FORTRAN IV. It uses an assembly language subroutine RDATA to read the data blocks produced by MR $\phi\phi$ 7. Figure 2 describes the actions of PUNCH.

After each logical record is written, a !EOD is written. After the last logical record, two !EODs are written (an EOF).

### B. Subroutine GETREC

GETREC is a FORTRAN subroutine which reads records produced by PUNCH into a data array. It tests sense switch 1 to determine the source device. The calling sequence is

```
CALL GETREC (IDAT(1), IDATL,999S)
```

when

IDAT(1) is the first element of the storage area. IDATL is the number of words to be read into the data storage area (currently 545). 999S is the address to which the routine will return when an end of all data (a double !EOD on the card reader, or an EOF on the DISK) is encountered.

If sense switch 1 is reset,

Data will be read from the Card Reader through F:8.

If sense switch 1 is set,

Data will be read from the DISK (DPEXEC) through F:7.

Figure 3 is a diagrammatic overview of GETREC.

## C. ACCTNG

ACCTNG is a FORTRAN program which produces a graphical summary of the execution data. The summary<sup>3</sup> consists of the following sections:

(I) Complete execution summary—All data are displayed.

(1) Instruction usage summary (Fig. 4)

The bar graph for the execution of an instruction shown in Fig. 4 is relative to the most used instruction. The bar graphs for indirect addressing (\*), indexing (X), and indexed indirect addressing (\*X) are relative to 100%. Unused instructions and addressing modes are not shown.

(2) Register usage summary

REGISTER USAGE refers to use of a register as an operand. INDEX REGISTER USAGE refers to use of a register for indexing. Bar graphs in each section are relative to the most used register in that section, as may be seen in Fig. 5.

(3) Execution classes

This is a one-page summary (see Fig. 6) of instruction usage. It is self-explanatory. Instructions are listed by their standard mnemonics.

(II) Condensed summary—Selected data are displayed

(1) Instruction usage summary

The condensed instruction usage summary is almost identical to the instruction usage summary. However, only instructions used more than 0.9% are listed, and the bar graphs for indirect, indexed, and indexed indirect addressing are relative to the individual instruction usage (see Fig. 7).

(2) Instruction overview

The instruction overview is a bar graph on a logarithmic scale of instruction usage. Instruc-

<sup>3</sup>The example used in this section is the execution of an optimized assembly language, FFT, written by Dr. Howard Rumsey. The FFT was executed twice in succession on a 256-point complex-element array.

tion mnemonics are on the bottom of the graph, and the corresponding usage is above. If an instruction was used at all, though not enough to appear on the graph, it has a “#” on the 1% grid. The total number of instructions is displayed on the 100% grid (see Fig. 8).

An instruction overview may be very revealing. Figure 8 shows an overview for an optimized assembly language FFT. Figure 9 shows an overview for an optimized FORTRAN IV FFT. Both were used to do a double transform on a 256-point complex-data array, and both were written by a highly skilled programmer. The FORTRAN program required almost twice as many instructions to do the transform as the assembly program. However, the absolute usage of floating point instructions is about the same (9–15% for assembly, 2 X 5–7% for FORTRAN).

ACCTNG will suppress all but the instruction overview if sense switch 2 is set. ACCTNG gets data through the subroutine GETREC (described in Section VB). Additionally, ACCTNG requires an index file to give it the op codes and mnemonics of the instructions, and control information for grouping the instructions together into logical blocks. The format of records in this file is as follows:

- (1) Op code, 4-letter mnemonic (blanks are significant)
- (2) Op code | sub code, 4-letter mnemonic
- (3) O, blanks

---

#### ABORT AND WARNING FORMAT:

HHHHHHHH

R $\phi$ ,RBLK $\phi$

RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR

R $\phi$ ,RBLK1

RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
RRRRRRRR	RRRRRRRR	RRRRRRRR	RRRRRRRR
LLLLLLLL	IIIIIIII	EEEEEEEE	CC $\phi\phi\phi\phi\phi\phi$

In this format,

Item (1) is the standard op code–mnemonic pair for an instruction. For example, the STORE REGISTER instruction, op code X'35', would be represented as “35,STW.”

Item (2) is for instructions which are actually a subclass of another instruction. In this case, the two-digit location of the data for the subcode is concatenated onto the two-digit op code. For example, the unconditional branch instruction (BCR,O) has data describing it stored in op code 01, and is represented as “6801,B.”

Item (3) establishes the end of a logical instruction group.

A diagrammatic summary of ACCTNG is shown in Fig. 10.

## VI. Error Messages and Other Miscellany

The following is a brief description of error messages and a possible explanation as to why they may occur, followed by a discussion on some conditions which have not been implemented and may cause program failure.

MR $\phi\phi$ 7 intercepts illegal and unimplemented instructions and gives a warning accordingly. MR $\phi\phi$ 7 also aborts on CALs that are the object of an EXU instruction, and flags EXU instructions which are nested. These events are accompanied by an error dump of the format shown below.

when

RBLK $\phi$  — Register block zero  $\rightarrow$  the program register block.  
RBLK1 — Register block one  $\rightarrow$  MR $\phi\phi$ 7 register block (or pseudo-register block).  
RRRRRRRR — Hexadecimal contents of the register.  
LLLLLLLL — Location of offending instruction.  
IIIIIII — Offending instruction.  
EEEEEEEE — SR3 from the last CAL 1 which aborted or errored.  
CC $\phi\phi\phi\phi\phi\phi$  — BYTE  $\phi$  contains the current condition code and floating control.

ERROR HEADERS (“HHHHHHHH” from above)

“ABRT CAL” — A CAL without an ERR/ABN exit has returned through the ERR/ABN exit with an abort level error code. — Abort.  
“BAD CAL” — An illegal CAL or a CAL with an illegal FPT code has been encountered. — Abort.  
“PRIV INS” — An attempt has been made to execute a privileged instruction. — Abort.  
“DATA INS” — An illegal/unimplemented instruction has been encountered. This includes CAL2 and CAL4 in BTM and CAL2, CAL3, CAL4 in BPM. — Abort.  
“EXU CAL” — An attempt to execute a CAL has been made. — Abort.  
“EXU EXU” — More than one level of EXU instruction. — Warning.

#### Miscellaneous Information

FPTs which are longer than 40 words will be truncated to 40 words. If there are any FPTs in the user program longer than 40 words, they should be reduced to 40 words or less.

SNAP CALs inserted by the monitor have their FPTs in 01 protection type. Since MR $\phi\phi$ 7 currently modifies the FPT, the monitor will abort on a memory protection violation when the SNAP is encountered. Therefore, !SNAP cards cannot be used in a region of the program to be executed by MR $\phi\phi$ 7. (M:SNAP FPTs are in 00 protection-type memory, and may be used.)

Since MR $\phi\phi$ 7 uses the M:EO DCB for saving data, this DCB must not be used by the program. Further, to insure

that the file (DEXEC) written through the M:EO DCB is not destroyed, the user must

- (1) exit to the monitor (through an M:EXIT, M:XXX, or M:ERR) while under the control of MR $\phi\phi$ 7

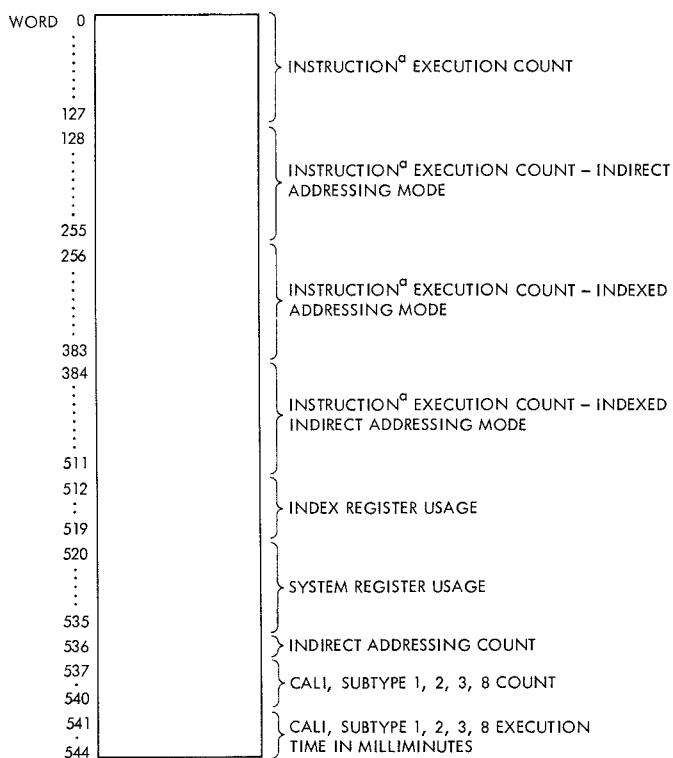
or

- (2) exit to the monitor via  
CALL WDATCLSE

or

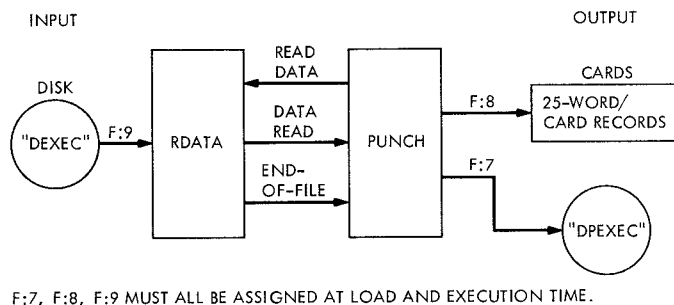
REF WDATCLSE  
B WDATCLSE

This closes the M:EO DCB with SAVE. All unclosed user output files are lost.

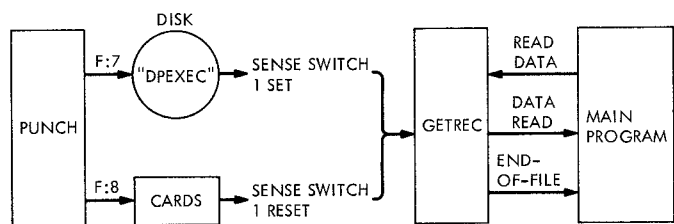


<sup>a</sup> AN ADDITIONAL INSTRUCTION SUBCLASS (BCR, O) HAS BEEN ASSIGNED OP CODE X'01'.

**Fig. 1. Data record format for DEXEC**



**Fig. 2. Overview of PUNCH**



**Fig. 3. Operational diagram of GETREC**



```

49=OR      .004%,      2!
48=ERR      .610%,      388!**
4B=AND      .007%,      4!
#####

25=S        .827%,      526!***
24=SF        .406%,      258!*
#####

3D=FAS      10.556%,      6717!*****
1D=FAL        .398%,      253!*
3C=FSS        9.952%,      6333!*****
3F=FMS      17.489%,      11129!*****
      * 4.58%,      510!***
1F=FML        .398%,      253!*
3E=FHS        .004%,      2!
#####

69=BFS        5.854%,      3725!*****
68=BCR        3.536%,      2250!*****
      * .71%,      16!
      X 23.29%,      524!*****
      *X .09%,      2!
01 B          .846%,      538!***
      * 2.97%,      16!**
      X 97.21%,      523!*****
      *X .37%,      2!
64=BDR        .002%,      1!
6A=BAL        .848%,      539!***
#####

04=CAL1      .005%,      3!
#####

CAL1,1:      .008MIN,      2!S *****
CAL1,2:      .000MIN,      0!S
CAL1,3:      .000MIN,      0!S
CAL1,8:      .000MIN,      0!S

```

Fig. 4. (contd)



REGISTER USAGE:

```

0 = .000%,      0:
:

1 = 1.877%,     1077:*****
:*****

2 = 8.987%,     5158:*****
:*****

3 = 1.795%,     1030:*****
:*****

4 = 7.999%,     4591:*****
:*****

5 = .004%,      2:
:

6 = .915%,      525:**
: **

7 = 1.784%,     1024:*****
:*****

8 = 25.990%,    14918:*****
:*****

9 = 9.556%,     5485:*****
:*****

A = 17.844%,    10242:*****
:*****

B = 7.152%,     4105:*****
:*****

C = 10.704%,    6144:*****
:*****

D = 3.581%,     2055:*****
:*****

E = .901%,      517:**
: **

F = .910%,      527:**
: **

```

INDEX REGISTER USAGE:

```

1 = 4.646%,     532:*****
:*****

2 = 35.816%,    4102:*****
:*****

3 = 35.851%,    4106:*****
:*****

4 = 1.197%,     137:**
: **

5 = .018%,      2:
:

6 = 4.593%,     526:*****
:*****

7 = 17.882%,    2048:*****
:*****

```

Fig. 5. Register usage summary format from ACCTNG

TOTAL INDIRECTS = 10796

TOTAL INSTRUCTIONS EXECUTED = 63636

RELATIVE USAGE OF INDIRECT ADDRESSING = 16.965 %

UNUSED INSTRUCTION SUMMARY:

LH	LCH	LAW	LFW	LCD	LAD	LS	LM
LCFY	LCF	XW	STH	STS	STM	STCF	ANLZ
INT	AW	SW	SD	MI	MH	DH	AWM
MTW	MTW	CB	CW	CD	CS	CLR	FSL
FDI	PSW	PLW	PRM	PLM	MSP	EXU	BIR
CALP	CAL3	CAL4					

INSTRUCTIONS USED LESS THAN 2.00%:

LI	LB	LAW	STB	SW	MW	DW	MTB
CI	CLM	BR	EAR	AND	S	SF	FAL
FML	FDS	R	BCR	BAL	CAL1		

INSTRUCTIONS USED MORE THAN .90%:

LI	LW	LD	STW	STD	AI	AW	AD
CW	FAS	FSS	FMS	BCS	BCR		

Fig. 6. Instruction class summary from ACCTNG



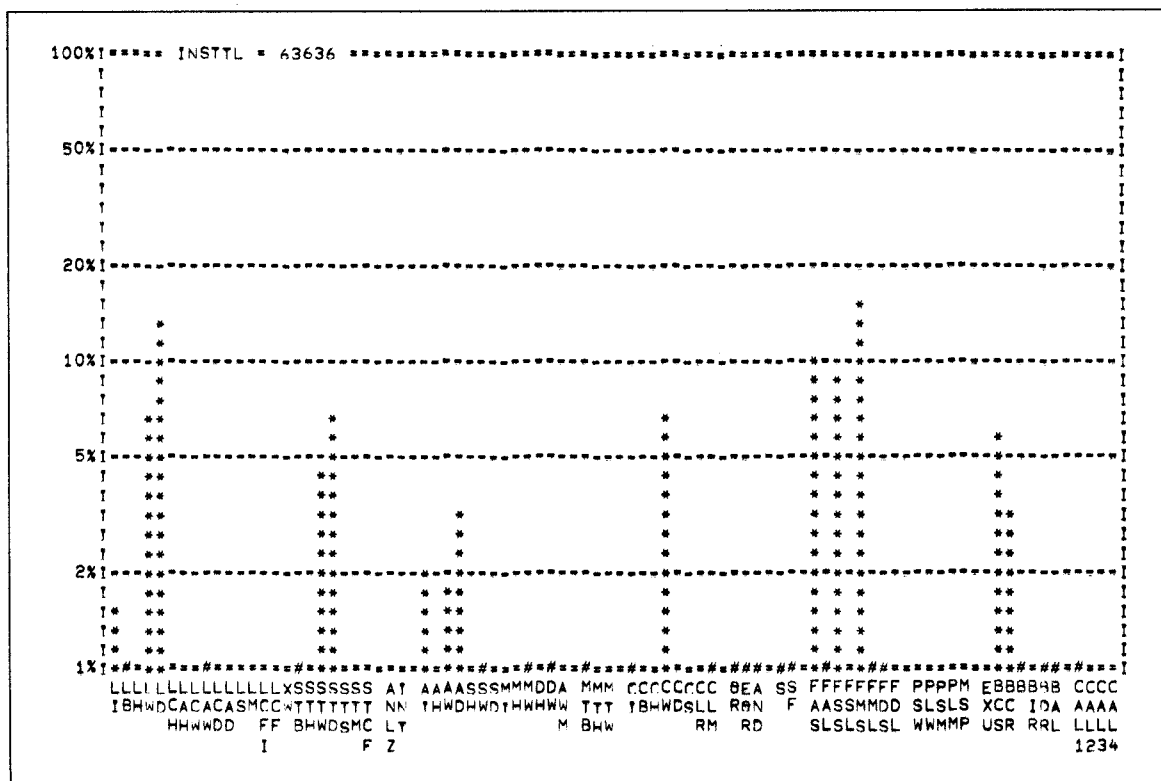


Fig. 8. Instruction overview of assembly-language FFT

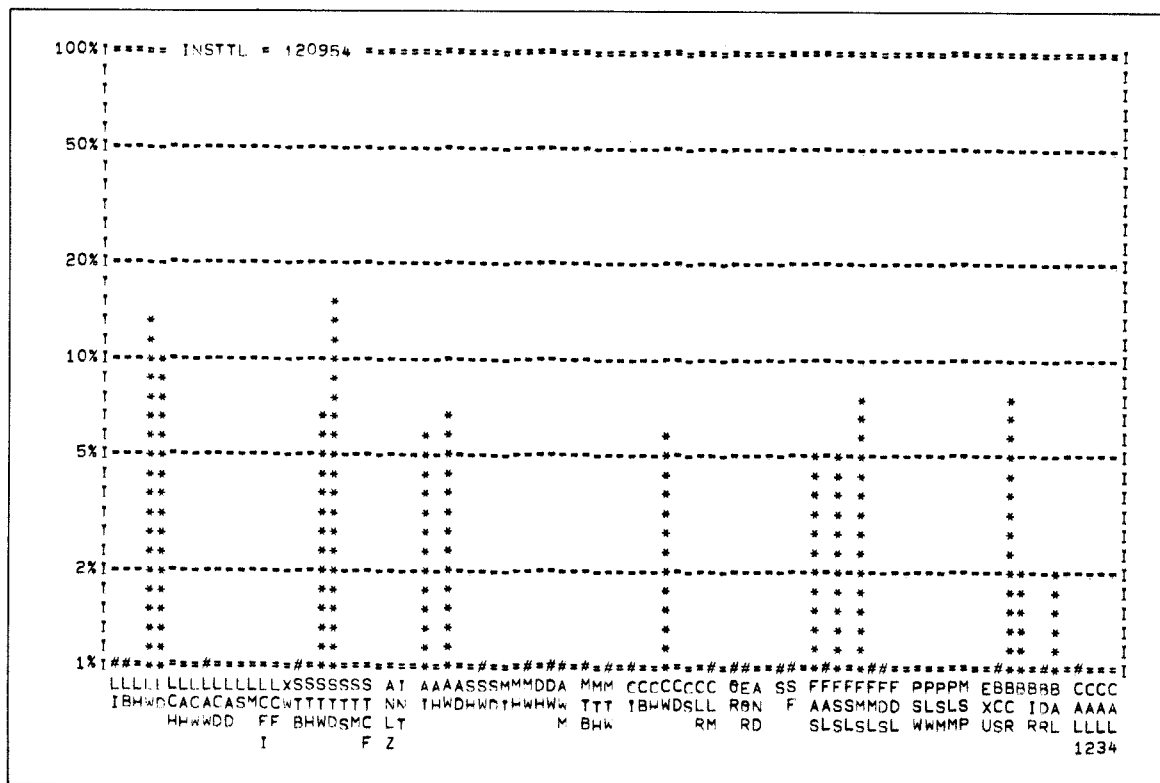
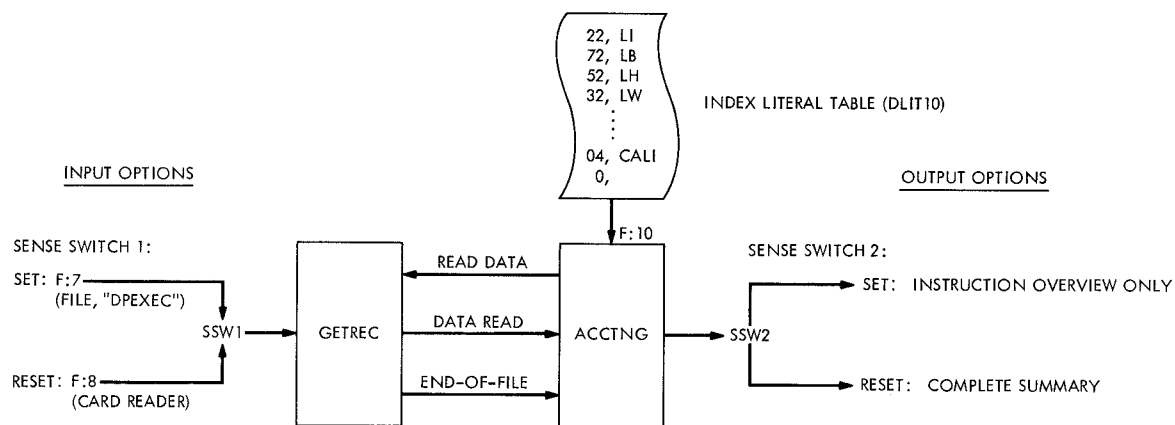


Fig. 9. Instruction overview of FORTRAN FFT



F:7, F:8, F:10 MUST BE ASSIGNED AT BOTH LOAD AND EXECUTION TIME.

Fig. 10. Diagrammatic summary of ACCTNG